

# 1 The Linear Least Squares Problem

**Introduction:** Determine unknown parameters  $c_1, \dots, c_n$  using  $m > n$  measurements with errors

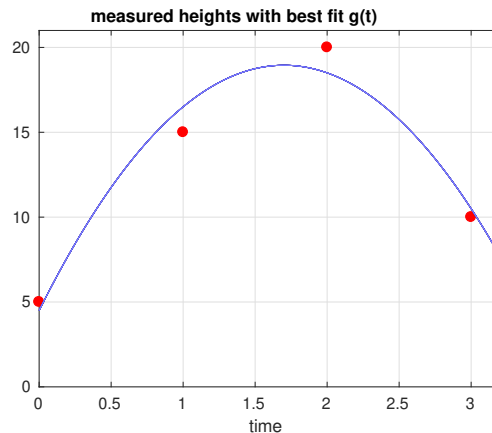
**Example:** We want to measure the acceleration  $g$  caused by gravity. If we throw an object into the air the height  $y$  is a quadratic function of time

$$y = g(t) = c_1 + c_2t + c_3t^2$$

with three unknown parameters  $c_1, c_2, c_3$  which we want to determine (then we obtain the acceleration from gravity as  $g = -2c_3$ ). In order to determine 3 unknown parameters we need at least 3 measurements. But we have measurement errors, so we want to perform a much larger number  $m$  of measurements.

We now perform our experiment: we throw an object into the air and measure its height at  $m$  different times  $t_1, \dots, t_m$ : We perform  $m = 4$  measurements and find the following data values (time  $t$  in seconds, height  $y$  in meters):

$t_j$	0	1	2	3
$y_j$	5	15	20	10



Note that there is no quadratic function  $g(t)$  which passes through all four points. We therefore want to find the function  $g(t) = c_1 + c_2t + c_3t^2$  which gives “the best fit” for the given data points.

In general the output value  $y$  depends on the input value  $x$  as follows:

$$y = g(t) \quad \text{with } g(t) = c_1g_1(t) + \dots + c_n g_n(t)$$

Here the functions  $g_1(t), \dots, g_n(t)$  are known, and we want to determine the unknown parameters  $c_1, \dots, c_n$ .

We perform  $m \geq n$  measurements and obtain data points  $(t_1, y_1), \dots, (t_m, y_m)$  where

$$y_j = g(t_j) + e_j, \quad j = 1, \dots, m$$

with **measurement errors**  $e_j$ . We assume that  $e_1, \dots, e_m$  are **small random errors** (we will be more precise below).

**Example:** If we want to fit measured points with a straight line, we have  $g(x) = c_1 \cdot 1 + c_2 \cdot t$  with  $g_1(x) = 1$  and  $g_2(x) = t$ . In this case we need  $m \geq 2$  points in order to be able to estimate  $c_1, c_2$ . Because of the random errors we should use  $m$  as large as possible. Note that for  $m > 2$  points we will not be able (in general) to find a straight line which passes through all the data points. We would like to find  $c_1, \dots, c_n$  which give the “best fit”.

For a certain choice  $c_1, \dots, c_n$  of the parameters we can measure the fit to the data values by the **residual vector**  $r = (r_1, \dots, r_m)^\top$  where

$$r_j := g(t_j) - y_j = c_1g_1(t_j) + \dots + c_n g_n(t_j) - y_j, \quad j = 1, \dots, m.$$

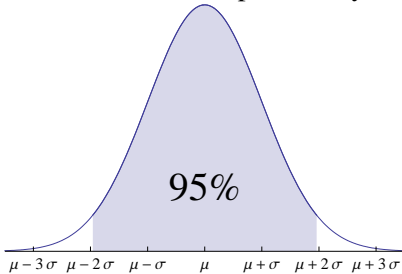
If the function  $g(t)$  were the true function, the observed values  $y_1, \dots, y_m$  would have errors  $r_1, \dots, r_m$ . Since large values of the errors are unlikely we want to pick  $c_1, \dots, c_n$  such that the residual vector  $r$  has a “small size”.

## Assumption about measurement errors

In order to understand which error vectors  $(e_1, \dots, e_m)$  are “likely” or “unlikely” we need to be more precise about the distribution of the errors. We first assume the following (in section 2 we will show that the least squares method also works under a weaker assumption)

**Assumption for errors  $e_1, \dots, e_m$ :**

- the errors  $e_1, \dots, e_m$  are independent
- the error  $e_j$  has **normal distribution** with a standard deviation  $\sigma$  (which is the *same* for all  $j = 1, \dots, m$ ): An error of size  $z$  occurs with a probability density of  $c e^{-z^2/(2\sigma^2)}$  which is the well-known bell-shaped curve:



Most errors observed in practice have (approximately) a normal distribution, since they are a sum of many independent sources (“central limit theorem”). The standard deviation  $\sigma$  describes how much the error is “spread out”. We have that  $|z| \leq 1.96\sigma$  with probability 95%. Note that large errors (“outliers”) are extremely unlikely.

For a certain choice  $c_1, \dots, c_n$  we obtain a residual vector  $(r_1, \dots, r_n)^\top$ . If  $(c_1, \dots, c_n)$  were the true values, the observed values  $y_1, \dots, y_m$  have a probability density which is given by the product of the individual densities:

$$C e^{-(r_1^2 + \dots + r_m^2)/(2\sigma^2)}$$

For the “maximal likelihood” (most plausible choice of  $c_1, \dots, c_n$ ) we should therefore minimize  $r_1^2 + \dots + r_m^2 = \|r\|_2^2$  (“least squares method”).

The least squares method may give bad results for  $c_1, \dots, c_n$  if our assumption is not satisfied. Two typical situations are:

- The errors  $e_i$  are normally distributed with known standard deviations  $\sigma_i$  which have different sizes (e.g., the measurement error is larger in certain intervals for  $x$ ). In this case we need to minimize

$$\frac{r_1^2}{\sigma_1^2} + \dots + \frac{r_m^2}{\sigma_m^2}$$

(“**weighted least squares method**”). We define  $\tilde{y}_j := y/\sigma_j$  and  $\tilde{a}_{jk} := a_{jk}/\sigma_j$ , then  $\frac{r_1^2}{\sigma_1^2} + \dots + \frac{r_m^2}{\sigma_m^2} = \|\tilde{A}c - \tilde{y}\|^2$  and we can use the normal algorithm (“ordinary least squares method”) with  $\tilde{A}$  and  $\tilde{y}$ .

- There are a few very large errors, so called “**outliers**”. This can be due to the fact that in addition to the standard error sources (small noise, measurement errors) there can be some rare large errors, e.g., if we accidentally knock against our delicate apparatus while performing our experiment. Since for the normal distribution large errors are extremely rare, outliers have a strong effect on the obtained parameters  $c_1, \dots, c_n$  and can spoil the result.

It turns out that the least squares method works well **even when the errors are not normally distributed**. We only need the following **properties for the errors**:

- the errors  $e_1, \dots, e_m$  are independent
- the mean (expectation) is zero:  $E[e_j] = 0$
- the variance is the same for  $j = 1, \dots, m$ :  $E[e_j^2] = \sigma^2$

## “Least squares method”

We want to find coefficients  $c_1, \dots, c_n$  such that the 2-norm  $\|r\|_2$  is minimal, i.e.,

$$F(c_1, \dots, c_n) := r_1^2 + \dots + r_m^2 = \text{minimal.} \quad (1)$$

Define the matrix  $A \in \mathbb{R}^{m \times n}$  by

$$A = \begin{bmatrix} g_1(t_1) & \cdots & g_n(t_1) \\ \vdots & & \vdots \\ g_1(t_m) & \cdots & g_n(t_m) \end{bmatrix}$$

then the residual vector is given by  $r = Ac - y$  and  $F(c_1, \dots, c_n) = \|Ac - y\|_2^2$ .

Therefore we can pose the least squares problem in the following form: Given a matrix  $A \in \mathbb{R}^{N \times n}$  and a right-hand side vector  $y \in \mathbb{R}^m$ , find a vector  $c \in \mathbb{R}^n$  such that

$$\|Ac - y\|_2 = \text{minimal.} \quad (2)$$

We will write  $\|\cdot\|$  for  $\|\cdot\|_2$  from now on.

## Normal Equations

Note that the function  $F(c_1, \dots, c_n)$  is a quadratic function of the coefficients  $c_1, \dots, c_n$ . Since this is a smooth function, at a minimum we must have that the partial derivatives satisfy

$$\frac{\partial F}{\partial c_1} = 0, \dots, \frac{\partial F}{\partial c_n} = 0. \quad (3)$$

Since  $F(c_1, \dots, c_n) = r_1^2 + \dots + r_m^2$  and  $r_j = a_{j1}c_1 + \dots + a_{jn}c_n - y_j$  we obtain using the chain rule

$$\frac{\partial F}{\partial c_1} = 2r_1 \frac{\partial r_1}{\partial c_1} + \dots + 2r_m \frac{\partial r_m}{\partial c_1} = 2r_1 a_{11} + \dots + 2r_m a_{m1} = 2[a_{11} \cdots a_{m1}] \begin{bmatrix} r_1 \\ \vdots \\ r_m \end{bmatrix} \stackrel{!}{=} 0$$

for the first equation in (3). All  $n$  equations in (3) together can therefore be written as

$$\begin{bmatrix} a_{11} & \cdots & a_{m1} \\ \vdots & & \vdots \\ a_{n1} & \cdots & a_{nn} \end{bmatrix} \begin{bmatrix} r_1 \\ \vdots \\ r_m \end{bmatrix} = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}, \quad \text{i.e., } A^\top r = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}.$$

These are the so-called **normal equations**. Since  $r = Ac - y$  we obtain  $A^\top(Ac - y) = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$  or

$$A^\top A c = A^\top y. \quad (4)$$

This leads to the following algorithm:

1. Let  $M := A^\top A$  and  $b := A^\top y$
2. Solve the  $n \times n$  linear system  $Mc = b$ .

Note that  $\|Ac - y\| = \min$  means that we want to approximate the vector  $y$  by a linear combination of the columns of the matrix  $A$ . If a column of  $A$  is a linear combination of some other columns, this column is “superfluous”, and leads to multiple solutions  $c$  which all give the same approximation  $Ac$ .

Therefore it makes sense to assume that the **columns of the matrix  $A$  are linearly independent**, i.e.,

$$Ac = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix} \implies c = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}. \quad (5)$$

This means that the rank of the matrix  $A$  is  $n$  (the rank is the number of linearly independent columns).

**Theorem 1.** *Assume that the columns of  $A$  are linearly independent. Then*

1. The normal equations have a unique solution  $c \in \mathbb{R}^n$ .
2. This vector  $c$  gives the unique minimum of the least squares problem: For  $\tilde{c} \in \mathbb{R}^n$  with  $\tilde{c} \neq c$  we have

$$\|A\tilde{c} - y\| > \|Ac - y\|$$

*Proof.* For (1.) we have to show that the matrix  $M = A^T A$  is nonsingular, i.e.,  $Mc = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix} \implies c = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$ . Therefore we

assume  $Mc = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$ . By multiplying with  $c^T$  from the left we obtain

$$\underbrace{c^T A^T}_{(Ac)^T} Ac = 0, \quad \text{i.e., } \|Ac\| = 0$$

which means  $Ac = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$ . Now our assumption (5) gives  $c = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$ .

For (2.) we let  $\tilde{c} = c + d$  with  $d \neq \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$  and have  $\tilde{r} := A\tilde{c} - y = A(c + d) - y = (Ac - y) + Ad = r + Ad$ , hence

$$\begin{aligned} \|\tilde{r}\|^2 &= (r + Ad)^T (r + Ad) = r^T r + 2(Ad)^T r + (Ad)^T (Ad) \\ &= \|r\|^2 + 2d^T \underbrace{A^T r}_{\begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}} + \underbrace{\|Ad\|^2}_{>0} > \|r\|^2 \end{aligned}$$

where  $A^T r = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$  by the normal equations, and  $\|Ad\| > 0$  because of  $d \neq \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$  and (5). □

The normal equations are the best method to solve a least squares problem by hand. On a computer with machine arithmetic it turns out that **sometimes this leads to a numerically unstable algorithm**. We can illustrate this by looking at the special case  $N = n$  with a square nonsingular matrix  $A$ . In this case  $\|r\| = \|Ac - y\|$  is minimized by solving the linear system  $Ac = y$ , and we have  $\|r\| = 0$ . Assume that  $A$  has a large condition number of about  $10^3$ , then typically  $A^T A$  has a condition number of about  $10^6$ . Therefore by solving the normal equations with matrix  $M = A^T A$  we will lose about 6 digits of accuracy. On the other hand we can just solve  $Ac = y$  and only lose about 3 digits of accuracy. Hence in this special case the algorithm with the normal equations is numerically unstable. It turns out that a similar loss of accuracy can also happen for  $N > n$  if we use the normal equations. Therefore we should use a different algorithm on a computer.

## Gram-Schmidt orthogonalization

Assume that we have a matrix  $A \in \mathbb{R}^{N \times n}$  with  $m \geq n$  and linearly independent columns, i.e.,  $\text{rank} A = n$ . We now want to find an **orthogonal basis**  $p^{(1)}, \dots, p^{(n)}$  of  $\text{range} A$ , i.e.,  $p^{(j)\top} p^{(k)} = 0$  for  $j \neq k$ .

We can do this in the following way: Denote the columns of  $A$  by  $a^{(j)}$ , i.e.,  $A = [a^{(1)}, \dots, a^{(n)}]$ .

Let  $p^{(1)} := a^{(1)}$ .

Let  $p^{(2)} := a^{(2)} - s_{12}p^{(1)}$  where we determine  $s_{12}$  such that  $p^{(2)} \perp p^{(1)}$ , i.e.,  $(a^{(2)} - s_{12}p^{(1)}) \cdot p^{(1)} = 0$  which yields

$$s_{12} := \frac{a^{(2)} \cdot p^{(1)}}{\|p^{(1)}\|^2}$$

We continue in the same way for  $p^{(3)}, \dots, p^{(n)}$ : We let

$$p^{(j)} := a^{(j)} - s_{1j}p^{(1)} - \dots - s_{j-1,j}p^{(j-1)} \quad (6)$$

where we determine  $s_{1j}$  such that  $p^{(j)} \perp p^{(1)}, \dots$ , determine  $s_{j-1,j}$  such that  $p^{(j)} \perp p^{(j-1)}$  yielding

$$s_{kj} := \frac{a^{(j)} \cdot p^{(k)}}{\|p^{(k)}\|^2} \quad k = 1, \dots, j-1 \quad (7)$$

By construction we have that the vectors  $p^{(1)}, \dots, p^{(n)}$  are orthogonal. By (6) we have

$$a^{(j)} = p^{(j)} + s_{1j}p^{(1)} + \dots + s_{j-1,j}p^{(j-1)}, \quad j = 1, \dots, n$$

which means in matrix form that

$$[a^{(1)}, \dots, a^{(n)}] = [p^{(1)}, \dots, p^{(n)}] \begin{bmatrix} 1 & s_{12} & \cdots \\ & \ddots & \ddots \\ & & 1 \end{bmatrix}$$

$$A = PS$$

where  $P = [p^{(1)}, \dots, p^{(n)}]$  has orthogonal columns and  $S$  is the upper triangular matrix with diagonal elements  $s_{jj} = 1$  and elements  $s_{kj}$  for  $j > k$  is given by (7).

## Solving the least squares problem using the decomposition $A = PS$

With the decomposition  $A = PS$  our problem becomes  $\|PSc - y\| = \min$ . With  $d = Sc$  we have  $\|Pd - y\| = \min$ , hence the normal equations give

$$P^\top Pd = P^\top y.$$

Since the matrix  $P^\top P$  is diagonal this is easy to solve: with  $b := P^\top y$  we obtain

$$d_j = \frac{b_j}{\|p^{(j)}\|^2} \quad j = 1, \dots, n$$

We can then find  $c$  by solving  $Sc = d$  using backsubstitution since  $S$  is upper triangular.

## Orthonormal basis and decomposition $A = QR$

We can normalize the vectors  $p^{(1)}, \dots, p^{(n)}$  to vectors  $q^{(1)}, \dots, q^{(n)}$  of length 1 by defining

$$q^{(j)} := p^{(j)} / \|p^{(j)}\|$$

Let  $Q := [q^{(1)}, \dots, q^{(n)}] \in \mathbb{R}^{m \times n}$ , then we have

$$A = QR \tag{8}$$

where  $R \in \mathbb{R}^{n \times n}$  is upper triangular with

$$R_{jj} = \|p^{(j)}\| > 0, \quad R_{kj} = \frac{a^{(j)} \cdot p^{(k)}}{\|p^{(k)}\|} \quad \text{for } j > k.$$

Since the columns of  $Q \in \mathbb{R}^{N \times n}$  are orthonormal we have  $Q^\top Q = I$ .

In **Matlab** we can compute the QR decomposition using **`[Q,R]=qr(A,0)`**

*Remark 1.*

1. The command **`[Q,R]=qr(A)`** (without “0”) gives a matrix  $Q \in \mathbb{R}^{m \times m}$  which has additional  $m - n$  columns so that all columns give an orthonormal basis of all of  $\mathbb{R}^N$ , and the matrix  $R \in \mathbb{R}^{N \times n}$  has  $N - n$  zero rows added at the bottom. We don’t need this for the least squares problem and use the “economy version” of the QR decomposition given by  $qr(A,0)$ .

Therefore the `qr` command allows us to **find an orthonormal basis for a subspace**  $W = \text{span}\{a^{(1)}, \dots, a^{(n)}\}$  of  $\mathbb{R}^N$  and **for its orthogonal complement**  $W^\perp$ : The first  $n$  columns of  $Q$  form an orthonormal basis of  $W$ , the remaining  $N - n$  columns form an orthonormal basis of  $W^\perp$ .

2. Actually Matlab does *not* use the Gram-Schmidt method for computing the QR decomposition since this may still be numerically unstable. Instead the so-called “Householder reflection method” is used (which I won’t explain here).

Therefore we obtain the following algorithm for solving the least squares problem:

1. Find a decomposition  $A = QR$  where  $Q \in \mathbb{R}^{m \times n}$  has orthonormal columns, and  $R$  is upper triangular.
2. Let  $b := Q^\top y$  and solve the upper triangular linear system

$$Rc = b$$

by back substitution.

In **Matlab** we can do this as follows:

```
[Q,R] = qr(A,0);  
c = R \ (Q' * y)
```

Matlab has a shortcut for this, we can just use the backslash operator:

```
c = A \ y
```