What's hat got to do with it? A hat trick of prisoner hat puzzles

Sam Lidz

April 2023

Adapted from slides by Eric Neyman

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Sam	1 10 2
Jann	LIUZ

What's hat got to do with it?

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Outline

The Finite Puzzle

- Statement
- Solution

2 Infinite Puzzle I

- Statement
- Detour: Equivalence Classes and the Axiom of Choice
- Solution

Infinite Puzzle II

- Statement
- Solution
- Boxes of Numbers
 - Statement
 - Solution

- Ten prisoners in a line, called P_1 through P_{10}
- Each prisoner is wearing a hat, which is either white or black
- P_i can see P_j 's hat whenever i < j
 - P_1 can see all hats but their own
 - P₁₀ can't see any hats
- The executioner asks each prisoner in order for their own hat color
- Prisoners who guess wrong are brutally executed

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If the prisoners strategize together beforehand, how many correct guesses can they guarantee?

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- If she sees an even number of black hats, she guesses white, otherwise she guesses black
- Then, P₂ knows whether there are the same parity of black hats in front of him as P₁ said, or different, and can thus correctly guess his own hat color.
- P_i knows the hat color of P₂, P₃,..., P_{i-1} from their guesses, and all the hats ahead of them by looking, so they can figure out their own hat color by the same parity comparison

Thus, every prisoner but P_1 is guaranteed to guess their own hat color correctly!

- Infinitely many prisoners in a line, called P_1, P_2, P_3, \ldots
- Each prisoner is wearing a hat, which is either white or black
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 - P_1 can see all hats but their own
 - P_{10} can see the hats of P_{11}, P_{12}, \ldots
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If the prisoners strategize together beforehand, what is the minimal number of incorrect guesses they can guarantee?

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Definition (equivalence relation)

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Example

Take our set to be $\mathbb{Z},$ and our equivalence relation to be congruence mod 3. Then 7 \sim 1, and $-1\not\sim$ 0.

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Remark

Equivalence classes partition the set on which the relation is defined into disjoint subsets, whose union is the original set.

The Axiom of Choice

If you have a collection of nonempty sets, there is a function that takes any set to something that's an element of it: $\forall C \text{ such that } S \in C \implies S \neq \emptyset, \exists F : C \rightarrow \bigcup C \text{ such that}$ $\forall S \in C, F(S) \in S$

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Remark

The Axiom of Choice lets us pick representatives from our equivalence classes, no matter how complicated these equivalence classes are.

Now we can finally solve the first infinite puzzle! (Remember the puzzle?) The following strategy will guarantee that at most one of the infinitely many prisoners will die!

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- Two sequences are equivalent if there are only finitely many differences between them

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 - Transitive? (a \sim b and b \sim c \implies a \sim c)
- For any pair of equivalent sequences, there's a last difference between them
- The equivalence relation gives equivalence classes, so using the axiom of choice the prisoners pick representatives

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- \bigcirc P_1 guesses white if she can see an even number of differences from the chosen representative, and black otherwise
- P₂ can now see if the parity of differences ahead of him is the same or different from the differences ahead of P₁, and can thus determine his hat color.
- P_i knows the hat color of P₂, P₃,..., P_{i-1} from their guesses, and all the hats ahead of them by looking, so they can figure out their own hat color by the same parity comparison

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The solution is exactly the same as the finite case! (Once we do a bunch of complicated set up)

- Infinitely many prisoners in a line, called P_1, P_2, P_3, \ldots
- Each prisoner is wearing a hat, which is either white or black
- P_i can see P_j 's hat whenever i < j
 - P₁ can see all hats but their own
 - P_{10} can see the hats of P_{11}, P_{12}, \ldots
- The executioner asks each prisoner *simultaneously* for their own hat color
- Prisoners who guess wrong are brutally executed

If the prisoners strategize together beforehand, what is the minimal number of incorrect guesses they can guarantee?

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Nobody has any information about their own hat, but by coordinating they can guarantee that "nearly every" prisoner survives!

- There is an infinite line of closed boxes, called B_1, B_2, B_3, \ldots
- In each box is a real number
- You may open any boxes, as long as you leave at least one closed
- After you're done looking at boxes, you must choose a closed box and guess the number inside it

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 - This is very similar to the equivalence relation on sequences of hats from earlier
- We can now pick representatives as before (again using the axiom of choice)
- Notice that for any pair of equivalent sequences, there's a last difference between them

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- Let *i* be greater than all the last differences of sequences we've seen, then open boxes B_{100(i+1)+1}, B_{100(i+2)+1},...

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- Let *i* be greater than all the last differences of sequences we've seen, then open boxes $B_{100(i+1)+1}, B_{100(i+2)+1}, \ldots$
- This gives us the equivalence class of the sequence in that row, so we guess that the number in box B_{100i+1} matches the representative

The chance that our chosen row had the latest last difference was 1%, so we guess correctly 99% of the time!

Questions?

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Infinite Puzzle II

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- 4 Boxes of Numbers
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